

On A Non-Archimedean Broyden Method

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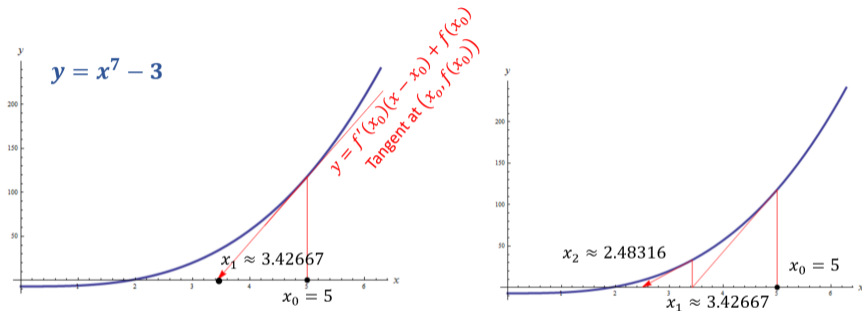
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Outline

1 introduction

2 The three contributions

Newton method over \mathbb{R} 

Approximation of a simple root x^* ($f'(x^*) \neq 0$).

Start from a guess x_0 , and compute $x_{n+1} = x_n - f'(x_n)^{-1}f(x_n)$.

System of m equations F in m unknowns: $x^* \in \mathbb{R}^m$, $F(x^*)$, $\text{Jac}_F(x^*)$ invertible

Start from a guess x_0 and compute: $x_{n+1} = x_n - \text{Jac}_F(x_n)^{-1}F(x_n)$ in \mathbb{R}^m .

Main result

- In 1965, Broyden suggested to replace the Jacobian matrix $\text{Jac}_F(x_n)$ by an approximation B_n of it → quasi-Newton method

$$x_{n+1} = x_n - B_n^{-1}F(x_n), \quad x_n, x_{n+1} \in \mathbb{R}^m, \quad B_n \in \text{Mat}_m(\mathbb{R}^m), \quad B_n \approx \text{Jac}_F(x_n)$$

- He chose $B_0 \approx \text{Jac}_F(x_0)$, then for B_1 a 1-dimensional deformation of B_0 , then for B_2 a 1-dim deformation of B_1 etc.
- 1965-today: numerous improvements, variants and generalizations. . .
- . . . but nothing for systems with coefficients in a complete valued field (p -adic or power series coefficients).

Main outcome : adapt the Broyden method to a system F with coefficients in a valued complete field → first non-archimedean quasi-Newton method

Broyden method in dim 1: secant method

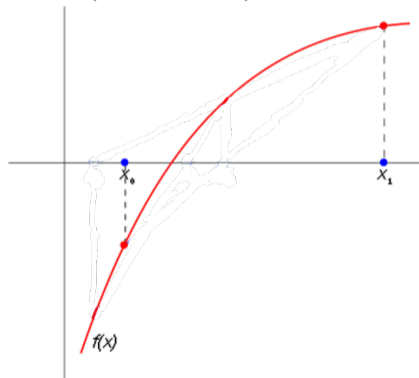
Broyden method generalizes the secant method (dimension 1) to dimension m .

Approximation of simple root x^*

$$f(x^*) = 0, \quad f'(x^*) \neq 0$$

Start with two guesses x_0, x_1 then,

$$x_{n+1} = x_n - \frac{x_n - x_{n-1}}{f(x_n) - f(x_{n-1})} f(x_n).$$



$$\text{Let } B_n^{-1} = \frac{x_n - x_{n-1}}{f(x_n) - f(x_{n-1})} \quad \rightarrow \quad \boxed{B_n(x_n - x_{n-1}) = f(x_n) - f(x_{n-1})}$$

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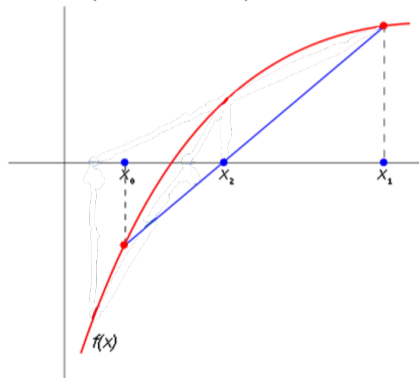
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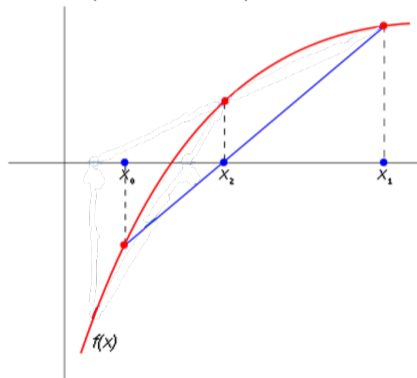
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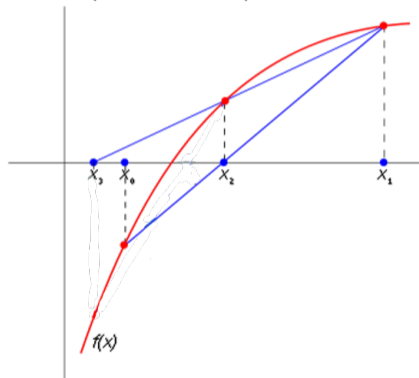
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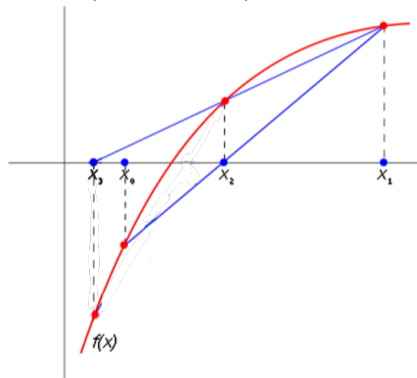
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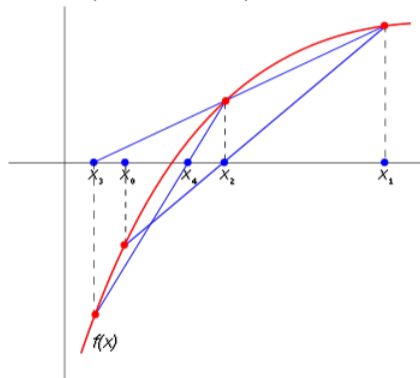
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Broyden's construction ($F : \Omega \subset \mathbb{R}^m \rightarrow \mathbb{R}^m$, $F(x^*) = 0$, $\text{Jac}_F(x^*) \neq 0$)

- Start with one guess x_0 and compute $B_0 \approx \text{Jac}_F(x_0)$.
- Then $x_{n+1} = x_n - B_n^{-1}F(x_n)$ with B_n verifying:
 - $B_n(x_n - x_{n-1}) = F(x_n) - F(x_{n-1})$ (Rewrite it $B_n s_{n-1} = y_{n-1}$)
 - $B_n z = B_{n-1} z$ for all z orthogonal to s_{n-1} .
- These two conditions suffice to characterize B_n from B_{n-1} , namely:

$$B_n = B_{n-1} + (y_{n-1} - B_{n-1}s_{n-1})u_{n-1}^T, \quad \text{where } u_{n-1}^T s_{n-1} = 1$$

- Broyden's choice : $u_{n-1} = s_{n-1} / \|s_{n-1}\|^2 = s_{n-1} / s_{n-1}^T \cdot s_{n-1}$.

Theorem 1 (Sherman-Morrison)

Computation of the inverse B_n^{-1} from B_{n-1}^{-1} can be done in $\approx 5m^2$

$$B_n^{-1} = B_{n-1}^{-1} - \frac{B_{n-1}^{-1}F(x_n)u_{n-1}^T B_{n-1}^{-1}}{u_{n-1}^T B_{n-1}^{-1}y_{n-1}}$$

Presentation of the Broyden method over \mathbb{R}

Start with $B_0 \approx \text{Jac}_F(x_0)$ for x_0 near to a non-singular solution x^* of F .

Broyden update (Iteration n)

$$\begin{aligned}x_{n+1} &= x_n - B_n^{-1}F(x_n), \\ B_{n+1}^{-1} &= B_n^{-1} - \frac{B_n^{-1}F(x_{n+1})u_n^T B_n^{-1}}{u_n^T B_n^{-1}y_n}, \quad (\text{vector } u_n \text{ verifying } u_n^T s_n = 1, \text{ like } u_n = \frac{s_n}{\|s_n\|^2})\end{aligned}$$

Cost of the n -th iteration (machine precision):

- $\approx 6m^2$ operations, and m evaluations of a scalar function (one evaluation of F)

Theorem 2 (Broyden — Dennis — Moré, 1973)

If F is sufficiently regular and x_0 sufficiently close to x^ then Broyden method produces*

a sequence that converges superlinearly to x^ :* $\lim_{n \rightarrow \infty} \frac{\|x_{n+1} - x^*\|}{\|x_n - x^*\|} = 0$.

Advantage of quasi-Newton methods over \mathbb{R}

Newton method:

- Computing the Jacobian matrix not always obvious... (1965)
- Needs to evaluate the m^2 entries of the Jacobian at each iteration.
- Needs to solve the system $\text{Jac}_F(x_n)s_n = -F(x_n)$ (in s_n) requires $O(m^3)$ (or $O(m^\omega)$ operations for m large).
- Quadratic convergence is fast, yet the region of quadratic convergence may be very small.

Whereas, with a Broyden update:

- Update costs $\approx 6m^2$ and one evaluation of F .
- \rightarrow no need to evaluate m^2 entries at each iteration.
- But has a slower rate convergence, that deteriorates with m .

However this drawback is mitigated with machine precision.

Outline

1 introduction

2 The three contributions

Adaptation to the non-archimedean setting (ex: p -adic or $k[[x]]$)

- With Broyden's original proposal: $B_{n+1} = B_n + \frac{(y_n - B_n s_n) s_n^T}{s_n^T \cdot s_n}$
($s_n = x_{n+1} - x_n$, $y_n = F(x_{n+1}) - F(x_n)$)
- requires a dot product which is an inner product over \mathbb{R}^m

- **Warning:** dot product is isotropic over \mathbb{Q}_p .
- However, if we use the previous formula:

$$B_{n+1}^{-1} = B_n^{-1} - \frac{B_n^{-1} F(x_{n+1}) u_n^T B_n^{-1}}{u_n^T B_n^{-1} y_n}, \quad u_n^T s_n = 1 \text{ we only need } u_n^T B_n^{-1} y_n \neq 0.$$

Contribution 1 (Adaptation)

Broyden update is adaptable over any complete valued field.

Idea: Write $s_n = (\sigma_1, \dots, \sigma_\ell, \dots, \sigma_m)^T$ and assume: $\text{val}(\sigma_\ell) = \min_{i=1, \dots, m} \text{val}(\sigma_i)$, ℓ is the **smallest** index verifying this property. Let $u_n = (0, \dots, \sigma_\ell^{-1}, \dots, 0)^T$. Then naturally $1 = u_n^T s_n$, and for technical reasons $u_n^T B_n^{-1} y_n \neq 0$.

Intermezzo: ultrametric norms

Let K be a complete discrete valued field:

- $K = \mathbb{Q}_p$, p -adic field or $K = k((X))$ field or Laurent series over a field k .
- Absolute value:
 - $|a| = p^{-\text{val}(a)}$ when $a \in \mathbb{Q}_p$, and $|a| = 2^{-\text{val}(a)}$ when $a \in k((X))$.
 - ultrametric inequality: $|x + y| \leq \max\{|x|; |y|\}$ (equality if $|x| \neq |y|$).
- Standard p -adic norms: $\|\vec{x}\| = \max\{|x_1|_p, \dots, |x_m|_p\}$.
 - Approximation x_n of a solution x^* : $\|F(x_n)\|$ is closer and closer to zero.

Lemma 3 (Operator norm on a matrix A)

$\|A\| = \max_{\|x\|=1} \|Ax\|$ is actually equal to the max-norm: $\|A\| = \max\{|\text{entries of } A|\}$.

- Over \mathbb{R} , matrices B_n minimize the Frobenius norm $\|B_{n+1} - B_n\|_F$
- \rightarrow Has been used to simplify the proof of superlinear convergence \rightarrow no such norm in the non-archimedean setting.

R-superlinear convergence

Contribution 2 (Convergence)

the non-archimedean Broyden method converges R-superlinearly of order at least

$\mu = 2^{1/2m}$:

$$\limsup_{k \rightarrow \infty} \|x_k - x^*\|^{1/\mu^k} < 1$$

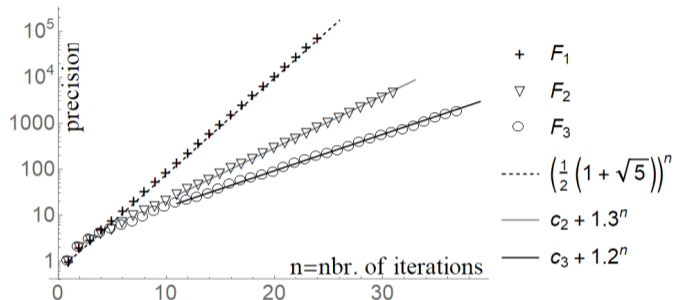
- **Meaning:** The sequence $\|x_k - x^*\|$ is not necessarily strictly decreasing. Essentially one can “think” that after $2m$ steps
 - $\|x_{n+2m} - x^*\| \leq C \|x_n - x^*\|^2$, for a constant C .
- **Experimental observations** suggest rather an “almost” Q-superlinear convergence of order $\alpha \approx 2^{1/m} \rightarrow \lim_{n \rightarrow \infty} \frac{\|x_{n+1} - x^*\|}{\|x_n - x^*\|^\alpha} \leq r$
- **Hypotheses:** $\text{Jac}_F(x^*)$ is invertible, and F verifies a kind of strong form of Taylor expansion at order 2 in a neighborhood U of x^* :
 - $\|F(y) - F(x) - \text{Jac}_F(x^*)(y - x)\| \leq c_0 \|y - x\|^2$, $\forall x, y \in U$, $c_0 > 0$
- **Remark:** This is the longest and most technical part of the article.

Experiments

F_1 : 2 polynomials, 2 unknowns

F_2 : 3 polynomials, 3 unknowns

F_3 : 4 polynomials, 4 unknowns



For $n = 2$, the order of superlinear convergence is $\Phi = \frac{1}{2}(1 + \sqrt{5})$: like in the secant method in one variable (archimedean setting or not).

- [E. Bach] Iterative root approximation in p-adic numerical analysis. *J. of Complexity* 2009

Implementation at finite precision: comparison with Newton

Here, the arithmetic cost depends on the precision. In case of Newton method (quadratic convergence):

$$\underbrace{x_{n+1}}_{2^{n+1} \text{ "digits"}} = \underbrace{x_n}_{2^n \text{ "digits"}} - \underbrace{\text{Jac}_F(x_n)^{-1} F(x_n)}_{\text{"digits" in the interval } [2^n, 2^{n+1}]}$$

- Therefore, no digits overlap in the sum \rightarrow easy to implement and analyze.
- \rightarrow The ratio (speed of convergence)/(precision gained) is somewhat optimal.
- Update of the Jacobian's inverse (by Newton method applied to $A \mapsto A^{-1}$)

$$\text{Jac}_F(x_{n+1})^{-1} = 2\text{Jac}_F(x_n)^{-1} - \text{Jac}_F(x_n)^{-1}\text{Jac}_F(x_{n+1})\text{Jac}_F(x_n)^{-1}$$

- Cost: $O(m^\omega)$ for matrix products.
 $O(mL)$ for the evaluation $\text{Jac}_F(x_{n+1})$ ([Baur-Strassen, 1980] \rightarrow if F is polynomial and can be valuated in L operations).

Implementation at finite precision: management of precision

Contribution 3 (Analysis of an implementation at finite precision)

Implementation, with polynomials as input, of Broyden method over \mathbb{Q}_p and $k((X))$ at finite precision and complexity analysis.

- Difficulty 1: The Broyden update uses a division \rightarrow bad for the precision

$$B_{n+1}^{-1} = B_n^{-1} - \frac{B_n^{-1}F(x_{n+1})u_n^T B_n^{-1}}{u_n^T B_n^{-1}y_n}, \quad x_{n+1} = x_n - B_n^{-1}F(x_n)$$

\rightarrow can be addressed by tracking p -adic intervals along these operations.

- Difficulty 2: convergence is not quadratic, and is not known.
 - Issue: If we extend precision of interval arithmetic too much, we loose efficiency.
 - however, the speed of convergence resembles to superlinear with an order of convergence α after a few iterations \rightarrow allows to guess without too much loss.

Complexity analysis - Comparison with Newton

- Notations:
 - $M(d)$ cost of multiplying two truncated power series/ p -adic integers in an interval of length d .
 - L number of arithmetic operations to evaluate the system F at any vector.
- n -th iteration of Newton method (interval of length 2^n): $O(M(2^n)(m^\omega + mL))$
- \rightarrow to reach precision $N \approx 2^\ell$, $O(M(N)(m^\omega + mL))$.

Assumption (order of superlinear convergence α for Broyden method)

Assume $\|x_{n+1} - x^*\| \leq r\|x_n - x^*\|^\alpha$ for $\alpha > 1$ and a constant $r > 0$.

- Cost of one iteration: $O(M(\alpha^n)(L + m^2)) \Rightarrow O(M(\frac{N}{\alpha-1})(m^2 + L))$.
- If we assume $\alpha \approx 2^{1/m}$ then $O(M(Nm)(m^2 + L))$ which is worse than Newton.

Application - Future work

Remark: Relaxed arithmetic [van der Hoeven et al.] (specific to Newton operator).

- p -adic and power series coefficients: quadratic convergence while maintaining $O(M(N)(m^2 + mL))$

Not better than Newton... so what's the point ?

- 1 Derivative-free: Jacobian is not easily accessible/computable \rightarrow divided-difference matrix.
- 2 Non-polynomial functions: [Baur-Strassen, 1980]'s theorem does not hold \rightarrow Jacobian is complicated to evaluate, may require up to m^2L operations to evaluate (instead of $O(mL)$).
- 3 Infinite dimensional problem: no Jacobian. Broyden's framework allows to work with finite dim. approximation [Kelley-Northrup,1988] [Kelley-Sachs, 1990] etc.